



Performance Analyses for Relaying LTE-A Connection Admission Control

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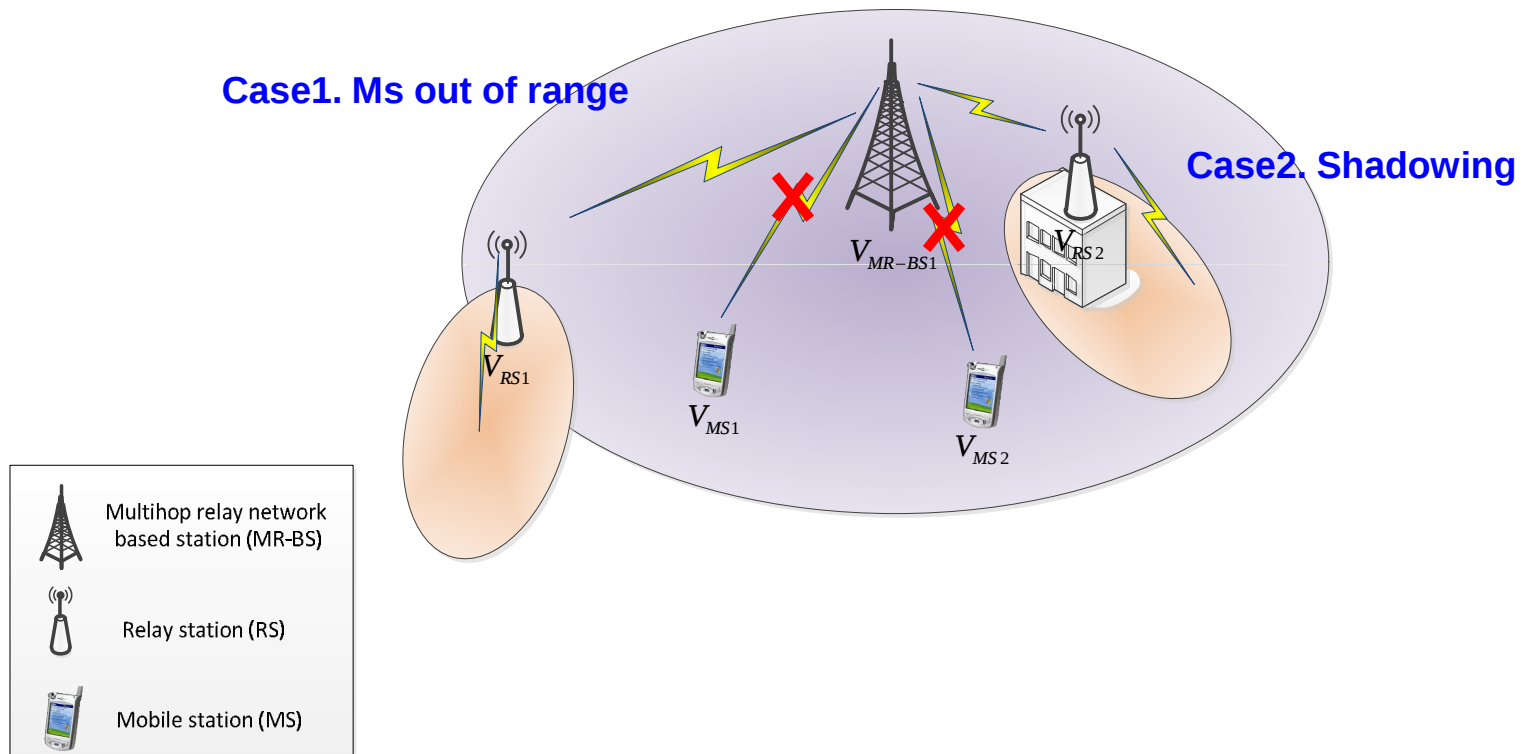
Outline

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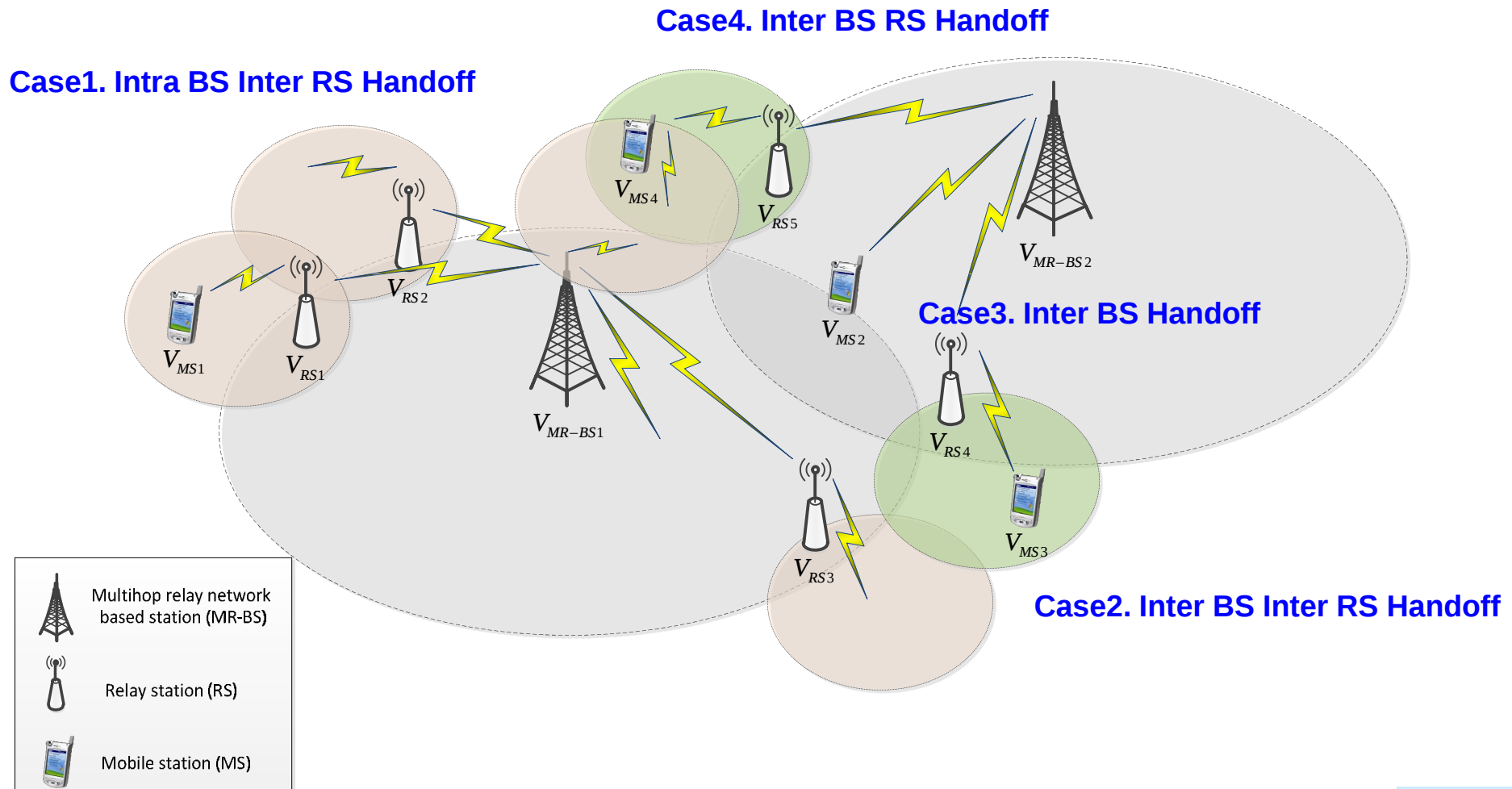


Overview





Handoff Cases in Relaying-based Cellular Networks





Issue

- Disadvantages of CAC in Relaying-based Cellular Networks (RCN)
 - 1. No efficient CAC in Relaying-based Cellular Networks
 - All of the handoff occur in both MR-BS and RS
 - 2. No adaptive CAC in Relaying-based Cellular Networks
 - Most CAC consider only the Ms type, and the user's bandwidth request or QoS





Related works

- Common Call Admission Control in wireless networks
 - **Complete sharing, Complete partition CAC**
 - [5] Salman A. AlQahtani and Ashraf S. Mahmoud, “**Dynamic radio resource allocation** for 3G and beyond mobile wireless networks,” *Computer Communications*, Vol. 30, Issue 1, pp. 41-51, Dec. 2006.
 - **Advantage**
 - Can achieve high Utilization
 - **Disadvantage**
 - Low system receipts
 - **Resource reserved CAC**
 - [7] M. Salamah and H. Lababidi “Dynamically adaptive **channel reservation scheme** for cellular networks,” *Computer Networks*, Vol. 49, Issue 6, pp. 787-796, Dec. 2005.
 - **Advantage**
 - Guarantee high priority user can use most resource
 - **Disadvantage**
 - It's hard to set the guard threshold





Related works (cont.)

– Cost based CAC

- **[14]** Wei shen and Qing-An Zeng, “**Cost-Function-Based** Network Selection Strategy in Integrated Wireless and Mobile Networks,” *IEEE Transactions on Vehicular Technology*, Vol. 57, Issue 6, pp.3778-3788, Nov. 2008.
- **Advantage**
 - Reducing blocking probability of ongoing call
- **Disadvantage**
 - The cost function weight was a static value





Related works (cont.)

- Analysis of Call Admission Control in wireless networks
 - **Analysis for Multi-service Networks**
 - [15] N. Lu and J. Bigham, “On utility-fair bandwidth adaptation for **multi-class traffic QoS** provisioning in wireless networks,” *Computer Networks*, Vol. 51, Issue 10, pp.2554-2564, July 2007.
 - [16] Yavuz E and V.C.M Leung, “Efficient Methods for Performance Evaluations of Call Admission control Schemes in **Multi-service** Cellular Networks,” *IEEE Transactions on Wireless Communications*, Vol. 7, Issue 9, pp. 3468-3476, Sep. 2008.
 - **Advantage**
 - Use multi-dimensional Markov chain to analysis multi-service class CAC in both CS and Guard channel
 - **Disadvantage**
 - It can't analysis for more than two traffic class networks





Motivation & Goals

- **Motivations**
 - Proposed a **Dynamic Cost-Reward based CAC** for MMR networks
 - Adaptive CAC for MR-BS and RS.
 - Decreasing the dropping probability of high priority users.
 - Maintain the system Utilization and more network reward.
- **Goals**
 - Using the **Dynamic Cost-Reward based CAC** to reach the goals as follow
 - 1. Low dropping probability
 - 2. Low FRL and GoS
 - 3. Maximum network reward





Proposed Approach

- **Proposed three steps to made decision**
 1. **Calculated BS cost**
 2. **Calculated inter arrival calls reward**
 3. **Compare Cost Reward to decision**
 - **Accept or reject**

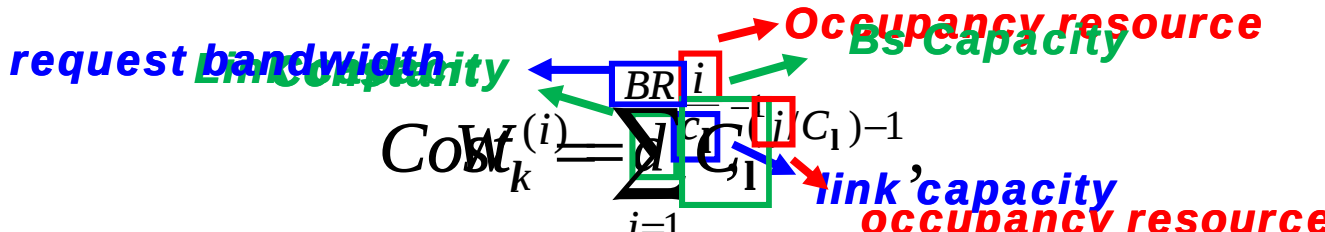




Step 1. Calculate BS cost

–[19] R. Gawlick, A. Kamath, S. Plotkin, and K. Ramakrishnan, “Routing and Admission Control in General Topology Networks,” *Technical report STAN-CS-TR-95-1548*, Stanford University, 1995.

- Competitive On-Line (COL)


$$Cost_k^{(i)} = \sum_{j=1}^N \frac{C_j}{C_1} (W_j(i) C_1)^{-1}$$
$$0 \leq W_l(i) \leq 1$$



Step 2. Calculate arrival calls reward

- Proposed a **Dynamic Reward Equation**

- Proposed **Dynamic Reward Equation**

- Motiv
- Goals

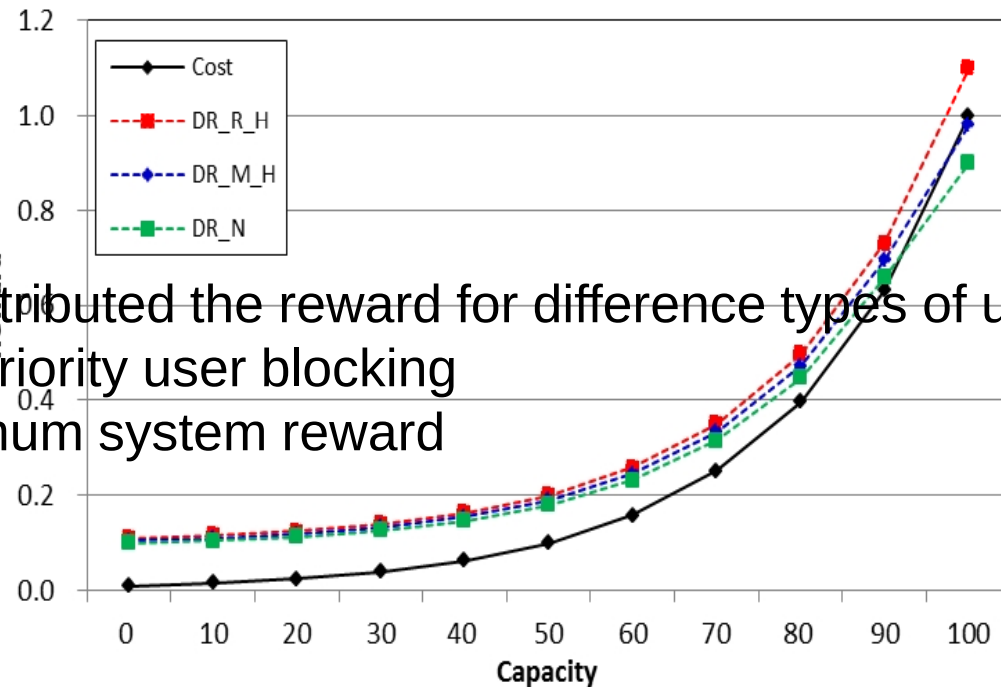
By distributed the reward for difference types of user to decrease high priority user blocking

— Maximum system reward

Connection ty

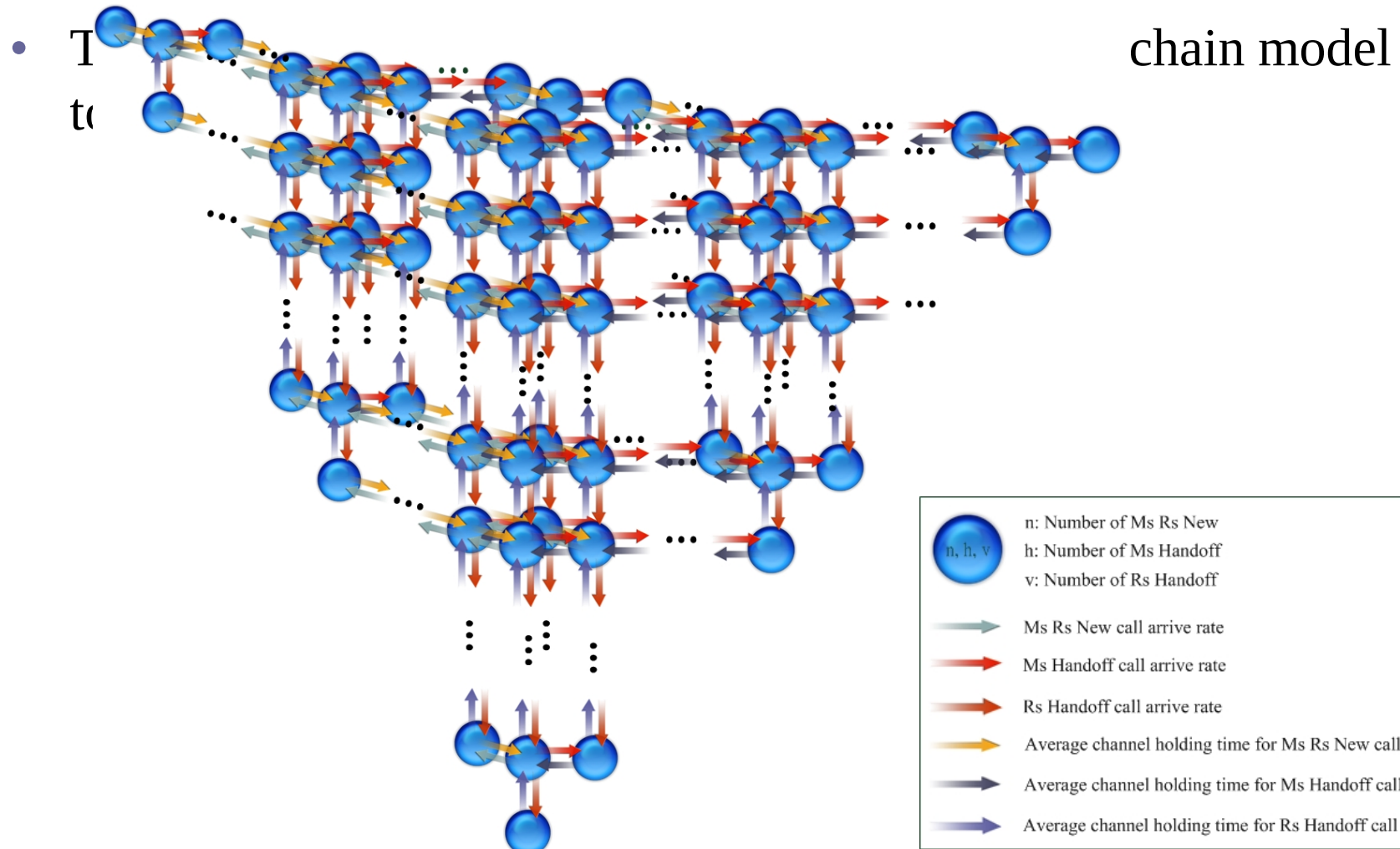
connection I in priority user

capacity





Markov Chain Analysis





Transition cases in CS and Guard

- The transition of all cases in the proposed Markov Chain model.

$$s(n, h, v ; n+1, h, v) = l_n^{Ms, Rs}$$

$$s(n, h, v ; n-1, h, v) = nm_n^{Ms, Rs}$$

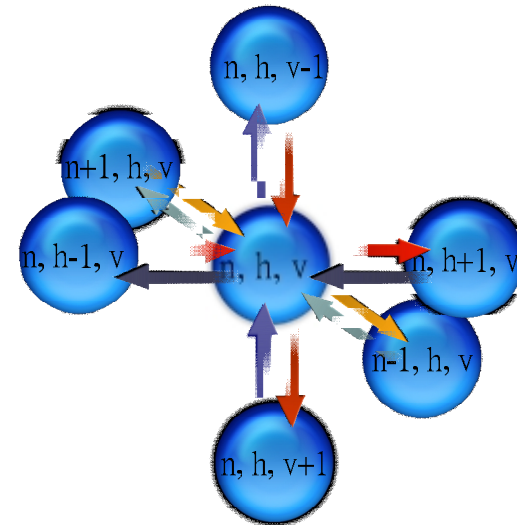
$$s(n, h, v ; n, h+1, v) = l_h^{Ms}$$

$$s(n, h, v ; n, h-1, v) = hm_h^{Ms}$$

$$s(n, h, v ; n, h, v+1) = l_h^{Rs}$$

$$s(n, h, v ; n, h, v-1) = vm_h^{Rs}$$

$$0 < n \leq t_n^{Ms, Rs}, 0 < h \leq t_h^{Ms}, 0 < v \leq t_h^{Rs}, n + h + v \leq Cap$$

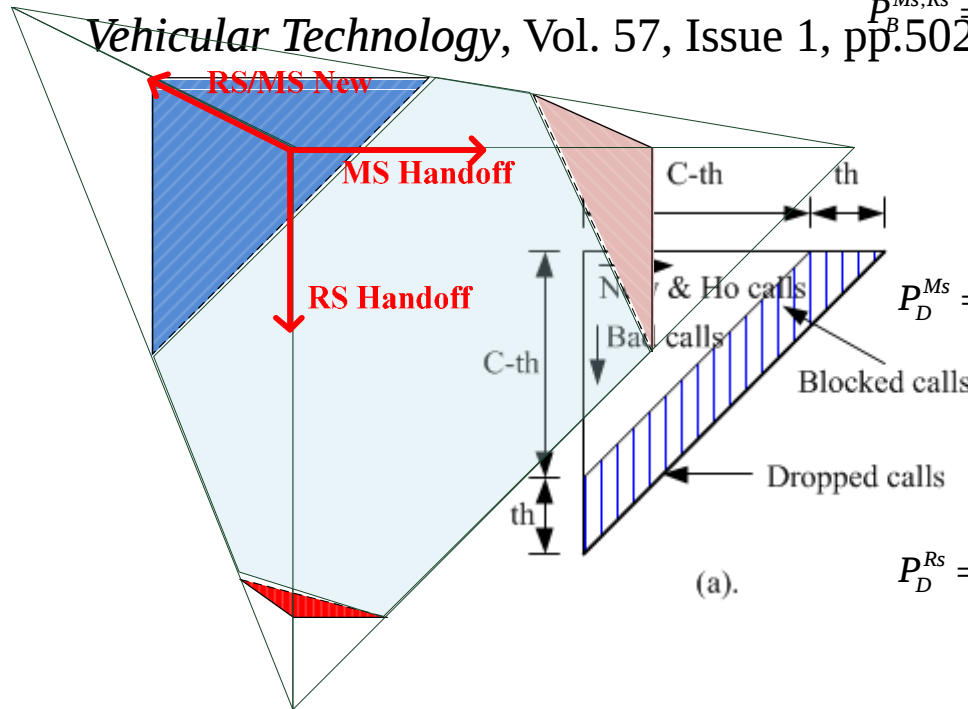




Performance in CS and Guard

- Blocking probability and Dropping probability can be calculated as follow.

- [22] B.B. Madan, S. Dharmaraja and K.S. Trivedi, "Combined Guard Channel and Mobile-Assisted Handoff for Cellular Networks," *IEEE Transactions on Vehicular Technology*, Vol. 57, Issue 1, pp.502-510, Jan. 2008.



$$P_D^{Ms, Rs} = \sum_{v=0}^{t_h^{Ms, Rs}} \sum_{n=0, 0 \leq Cap-v-n \leq t_h^{Ms, Rs}} \frac{r_v^v r_n^n r_h^{Cap-v-n}}{t_h^{Ms, Rs} v! n!} P(0,0,0) + \sum_{v=0}^{t_h^{Ms, Rs}} \sum_{h=0}^{t_h^{Ms, Rs}} \frac{r_v^v r_h^h r_n^{Cap-v-h}}{t_h^{Ms, Rs} v! h! n!} P(0,0,0)$$

$$P_D^{Ms} = \sum_{v=0}^{t_h^{Ms}} \sum_{n=0, 0 \leq Cap-v-n \leq t_h^{Ms}} \frac{r_v^v r_n^n r_h^{Cap-v-n}}{v! n! (Cap-v-n)!} P(0,0,0) + \sum_{v=0}^{Cap-t_h^{Ms}-1} \sum_{n=0}^{Cap-t_h^{Ms}-v-1} \frac{r_h^{t_h^{Ms}} r_v^v r_n^n}{t_h^{Ms} v! n!} P(0,0,0)$$

$$P_D^{Rs} = \sum_{v=0}^{t_h^{Rs}} \sum_{h=0, 0 \leq Cap-v-h \leq t_h^{Rs}} \frac{r_v^v r_h^h r_n^{Cap-v-h}}{v! h! (Cap-v-h)!} P(0,0,0) + \sum_{h=0}^{Cap-t_h^{Rs}-1} \sum_{n=0}^{Cap-t_h^{Rs}-h-1} \frac{r_v^{t_h^{Rs}} r_h^h r_n^n}{t_h^{Rs} h! n!} P(0,0,0)$$





Original State $P(0,0,0)$

- Obtained above equations, the original state can be presented by the following equations :

$$P(0,0,0) = \left(\sum_{n=0}^{t_n} \sum_{h=0}^{t_h^{Ms}-n} \sum_{v=0}^{t_h^{Rs}-n-h} \frac{r_n^n}{n!} \frac{r_h^h}{h!} \frac{r_v^v}{v!} \right)^{-1},$$

$$0 < t_n^{Ms, Rs} < t_h^{Ms} < t_h^{Rs} \leq Cap, \quad n + h + v \leq Cap, \quad r_i = \frac{l_i}{m_i}$$





Transition cases in Cost-Reward based CAC

- The transition of all cases in the proposed Cost-Reward based CAC.
- Not only satisfy by transition in CS and Guard, but also satisfy by the Cost Reward.

$$r_{n+1,h,v} \geq C_{n+1+h+v}$$

$$s(n, h, v ; n+1, h, v) = l_n^{Ms, Rs}$$

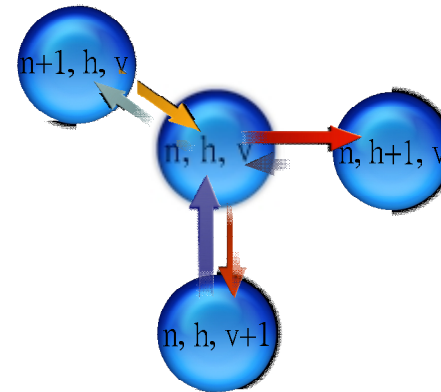
$$r_{n,h+1,v} \geq C_{n+h+1+v}$$

$$s(n, h, v ; n, h+1, v) = l_h^{Ms}$$

$$r_{n,h,v+1} \geq C_{n+h+v+1}$$

$$s(n, h, v ; n, h, v+1) = l_h^{Rs}$$

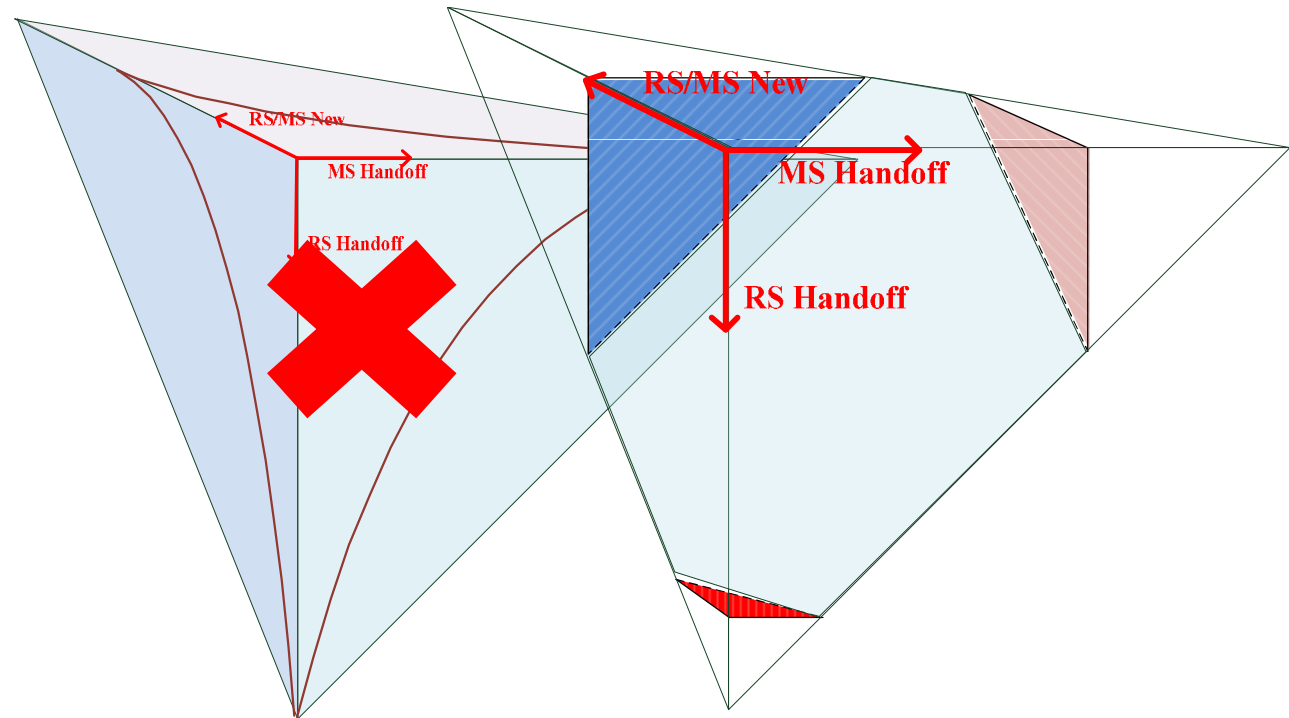
$$0 < n \leq t_n^{Ms, Rs}, 0 < h \leq t_h^{Ms}, 0 < v \leq t_h^{Rs}, n + h + v \leq Cap$$





Performance ***in Cost-Reward based CAC***

- Blocking probability and Dropping probability can not be calculated.





Performance in Cost-Reward based CAC (cont.)

–**[23]** V.B. Iversen et al. “Teletraffic Engineering. Handbook,”
ITU–D SG 2/16 & ITC, Jan. 2005.

- State probability can be calculated as follow.

$$P(i, j) = P(i) * P(j) = Q \frac{A_1^i}{i!} \frac{A_2^j}{j!}$$

- and

$$Q(n) = \sum_{i=0}^n q(i),$$

$$q(x) = \frac{1}{x} \sum_{j=1}^N A_j \cdot q(x-1), \quad q(0) = 1,$$



Performance in Cost-Reward based CAC (cont.)

- State probability:

$$p(x) = \frac{q(x)}{Q(n)}, \quad 0 \leq x \leq n,$$

- We modify the Equation to appropriate the model:

$$p(n, h, v) = \frac{q(n, h, v)}{Q(n, h, v)},$$

$$0 < n \leq t_n^{Ms, Rs}, \quad 0 < h \leq t_h^{Ms}, \quad 0 < v \leq t_v^{Rs}, \quad n + h + v \leq Cap,$$



Numerical Results

- The proposed approach will compare with these approaches as follow
 - 1. Complete Sharing CAC(**CS**)
 - 2. Resource reserved CAC(**Guard_th**)
 - 3. Complete Partition CAC(**CP**)
 - 4. Static Reward CAC (**SR**)
- Compare Performance Metrics as follow
 - 1. **FRL**
 - 2. **Grade of Service**
 - 3. **Reward**
 - 4. **Utilization**





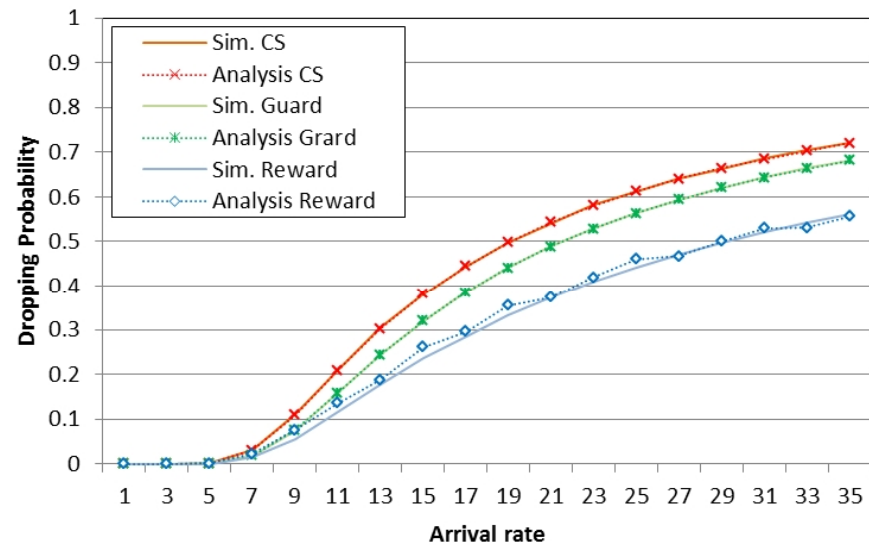
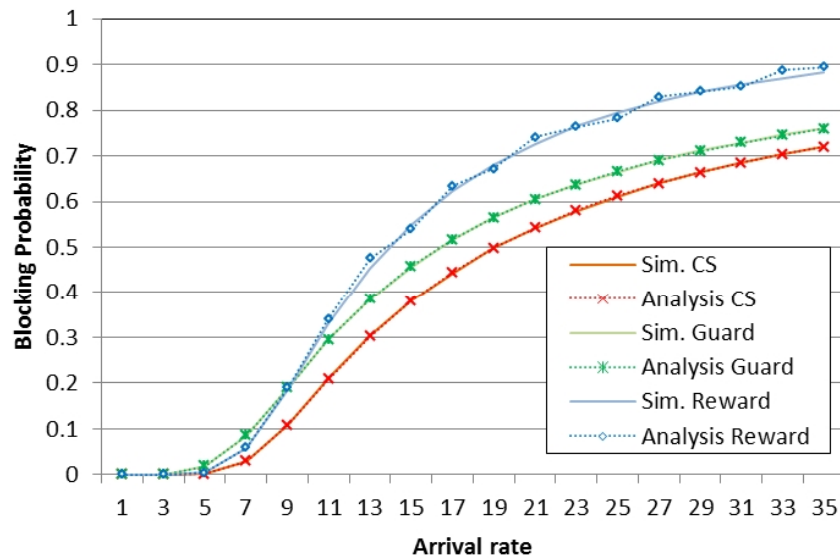
Simulation & Analysis Parameters

<i>Parameter</i>	<i>Value</i>
Simulation time (Second)	1000
BS network capacity	40 Mbps
Number of Mobile nodes (M_NDS)	20, 40
Number of Mobile nodes (M_NDS) l	10
Arrival rate of Poisson distribution ,	3-18
Average holding ^m time of Exponential distribution,	1
Guard threshold	0.7
CP threshold (High, Low)	0.6, 0.4





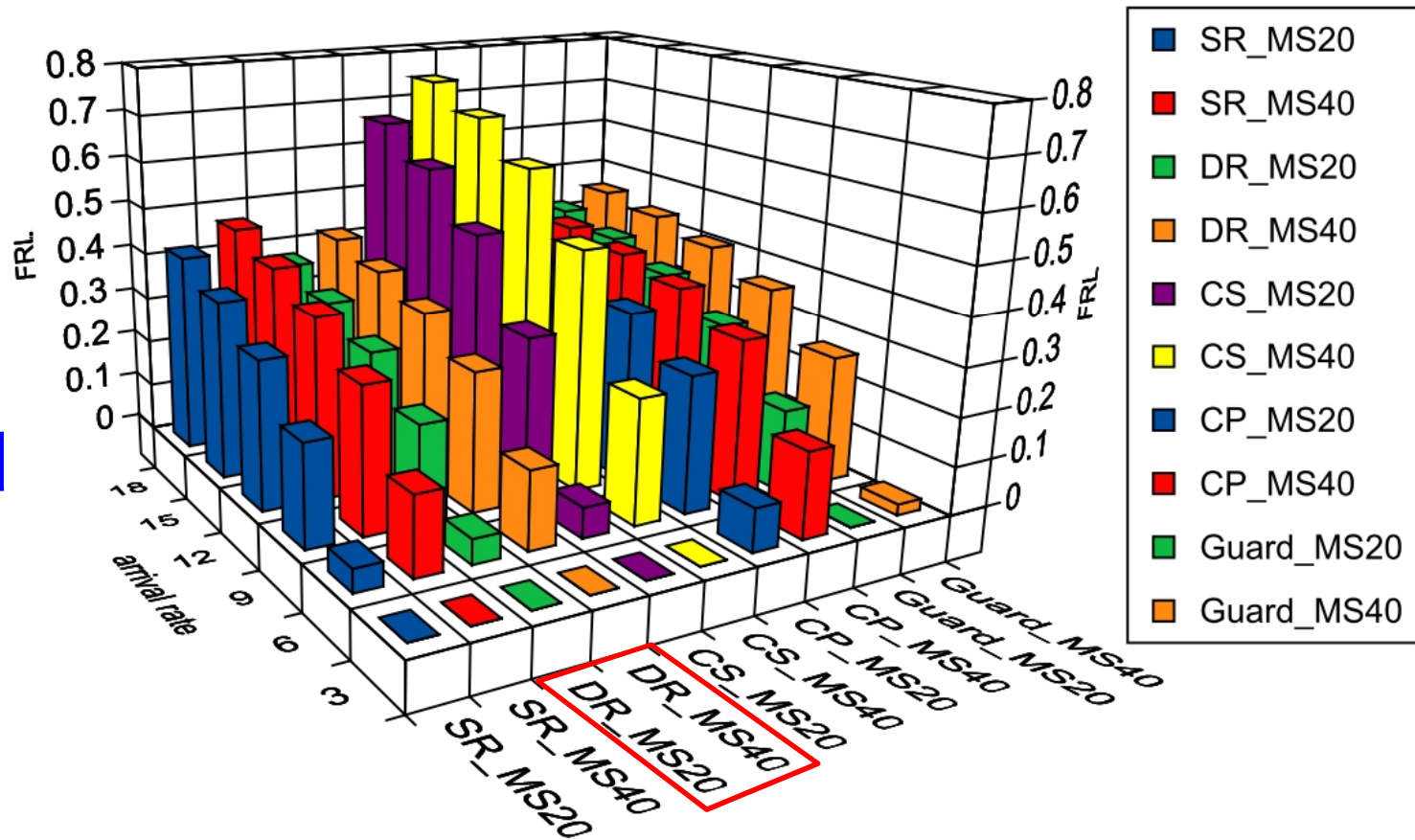
Blocking / Dropping Probability – Simulation & Analysis





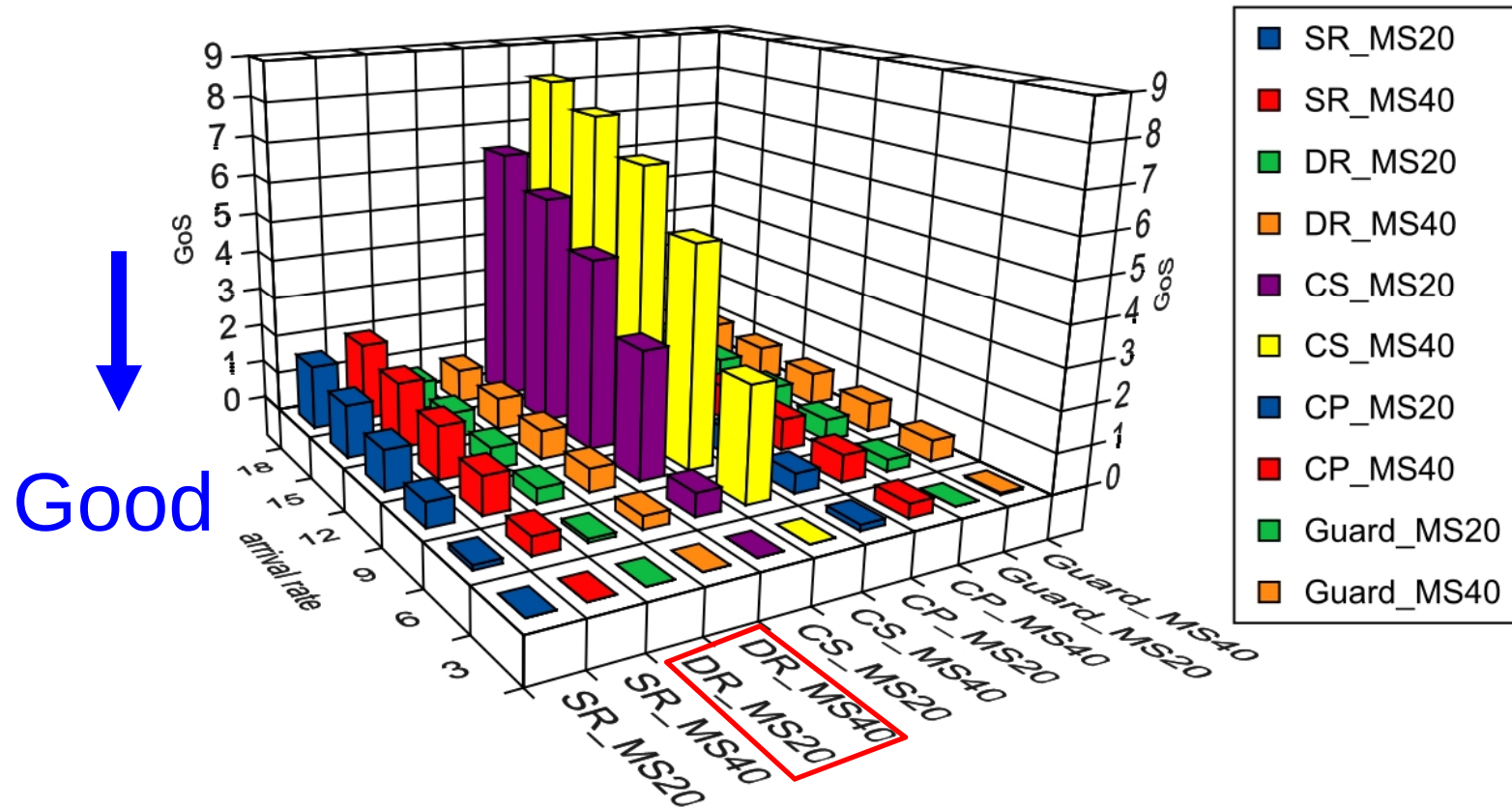
Performance – FRL

Good



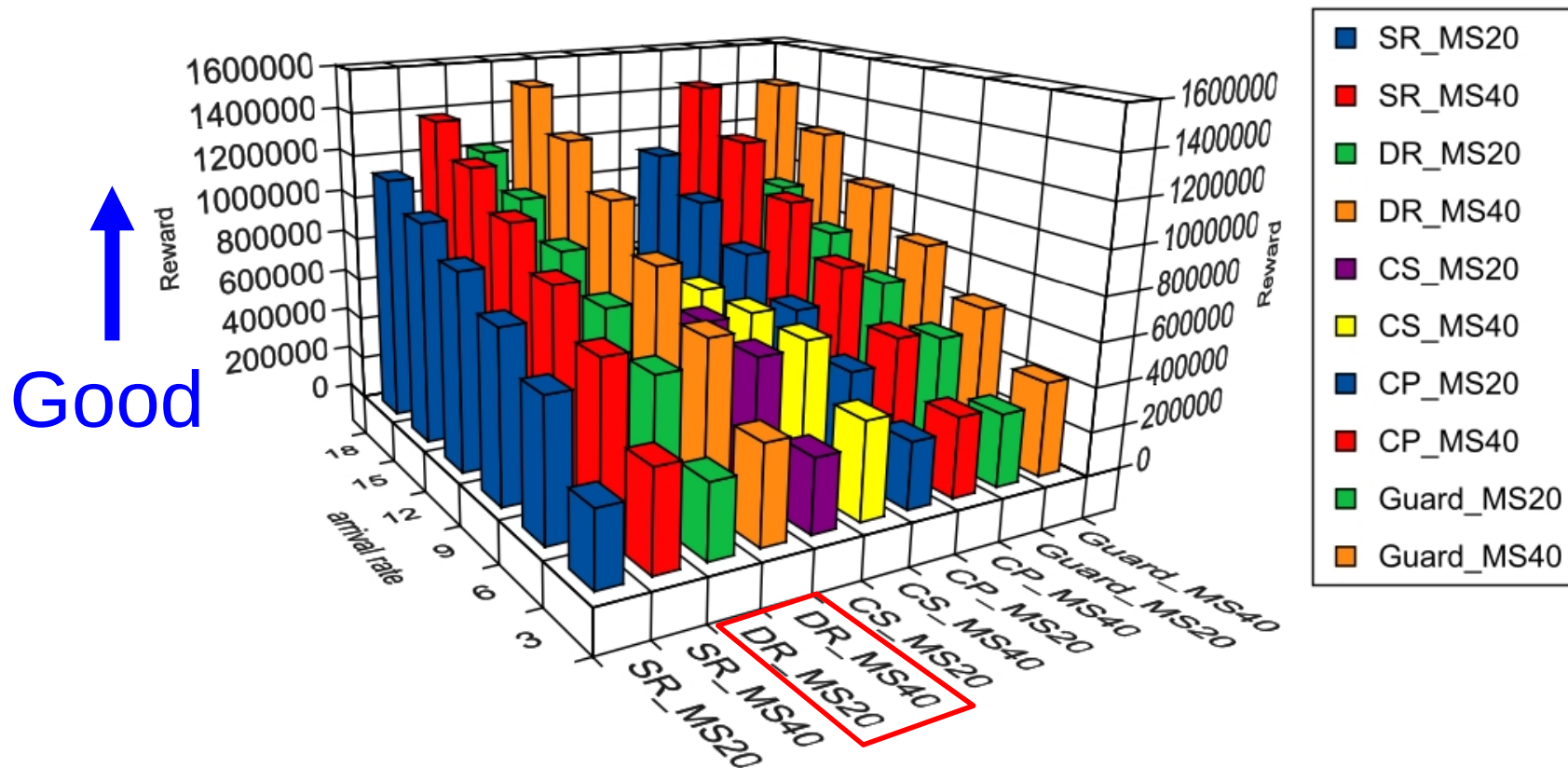


Performance – GoS



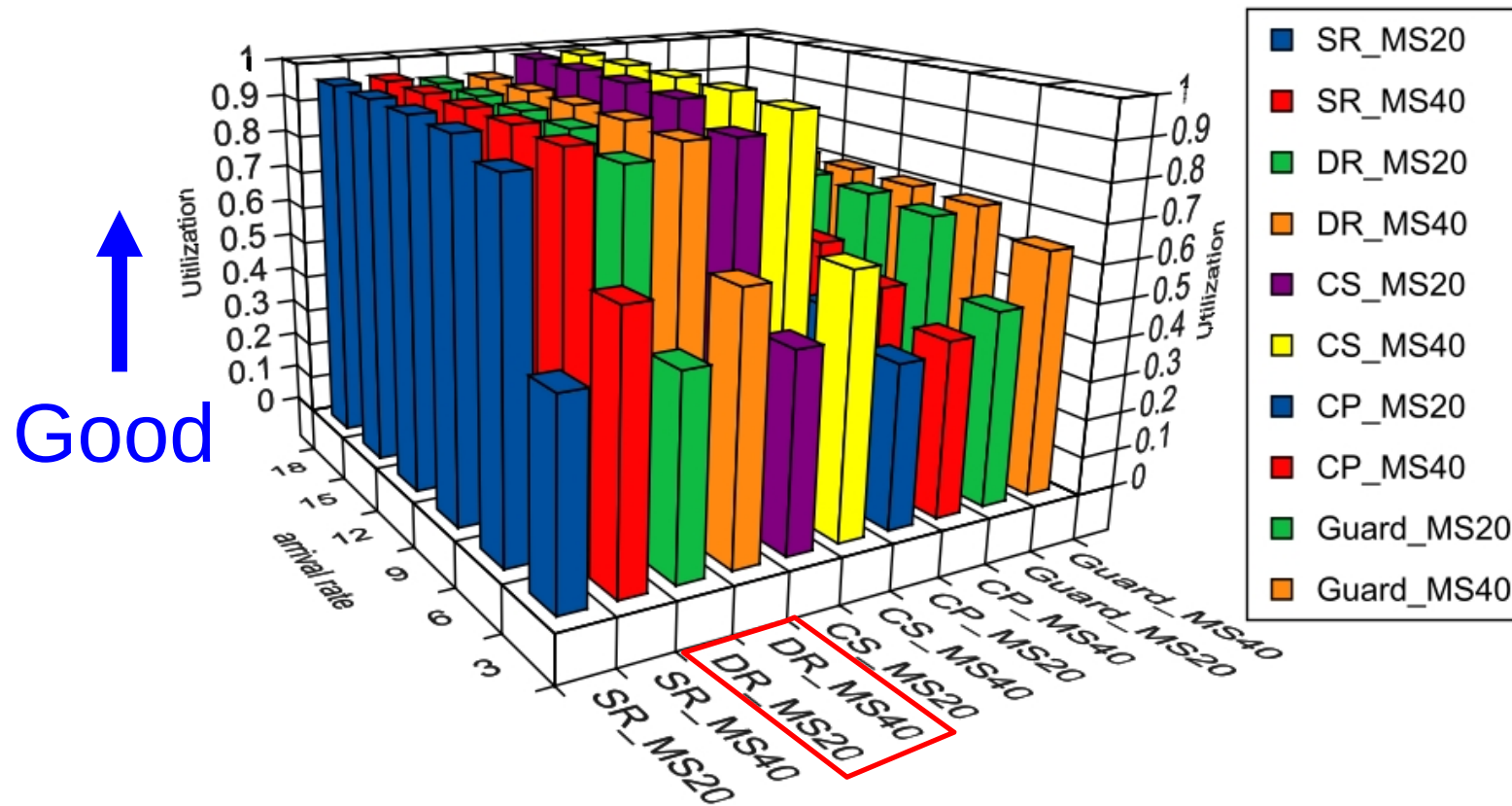


Performance – Reward





Performance – Utilization





Conclusions

- Step 1. Calculate Bs Cost
- Step 2. Calculate arrival call Reward
- Step 3. Compare Cost Reward to decision

- Reaches goals :
 - 1. **Dropping probability** (Dynamic cost-reward based CAC)
 - 2. **Reward** (Dynamic cost-reward based CAC)
 - 3. **Performance analysis** (Multi-dimensional Markov chain model)

- This research **not only** simulation the network condition **but also** analysis the currently of performance metrics

- The numerical results to show the requests of simulation is **closed** to mathematical analysis





Q & A ***Thanks***

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